

Comparison of TCP Performance Models

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Motivation: Internet on Air (IoA)

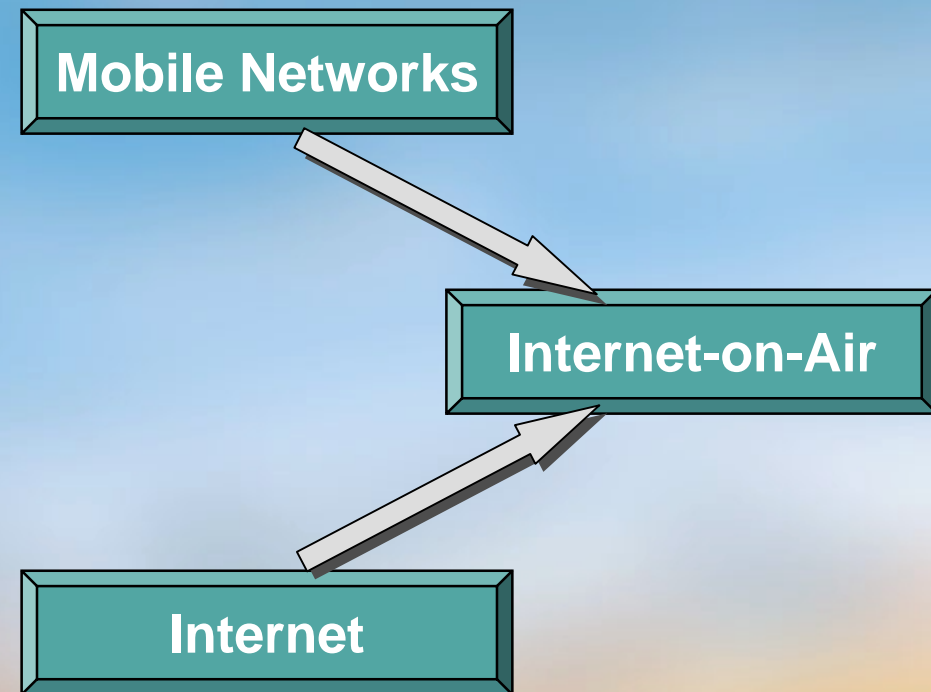
IoA: The Evolution Target of Mobile Networks and Internet

Trends in Mobile Networks:

- ≡ IP transport in the backbone; IP in RAN
- ≡ Transport voice & data over IP
- ≡ Terminate IP in the mobile host
- ≡ Separation: Transport \leftrightarrow Control

Trends in the Internet:

- ≡ Enable wireless access, support mobility
- ≡ QoS beyond “Best Effort”
- ≡ Security and AAA



TCP traffic dominant in the Internet \rightarrow Wireless TCP important in IoA

Challenge: Traffic/Performance Models for TCP

[Even without wireless links]

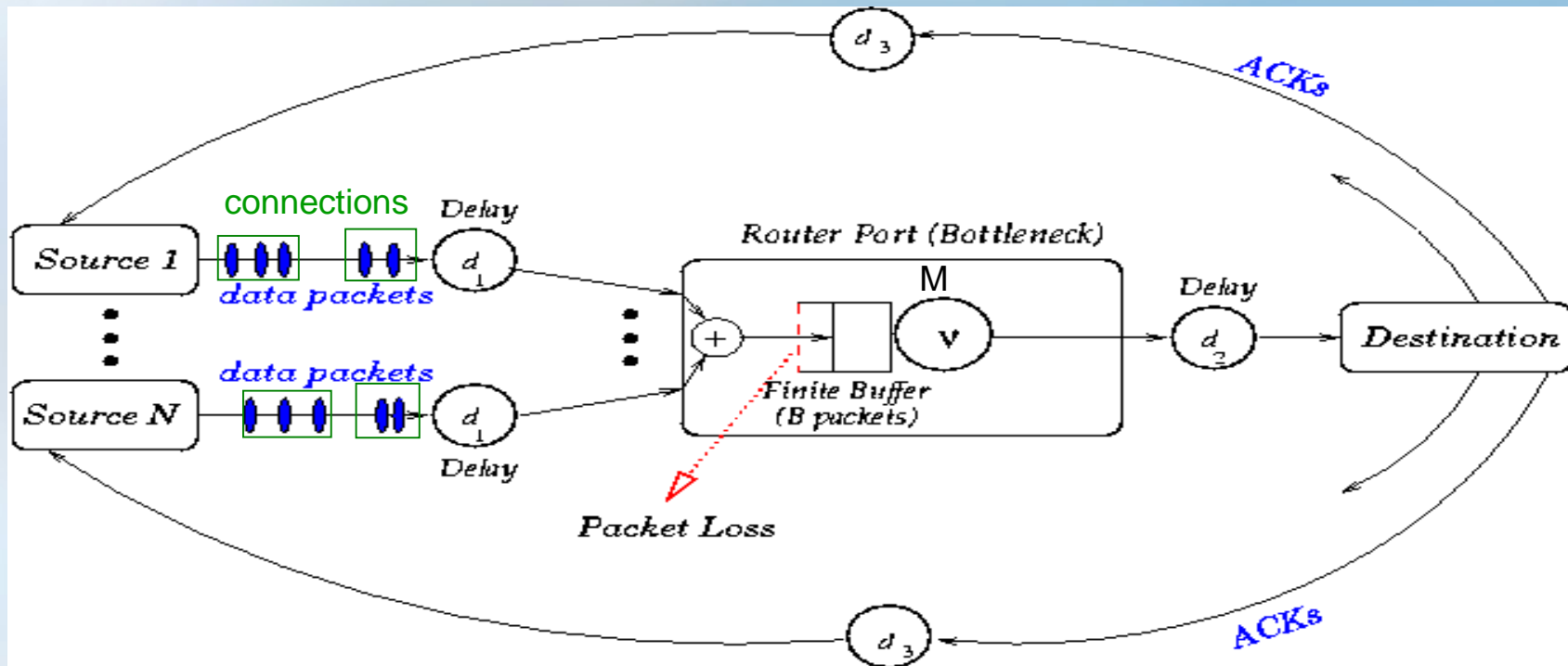
End-to-End flow-control (using Ack packets)

- Feedback: network behavior (congestion) \leftrightarrow ingress traffic
- No separation traffic model/network model possible
- New traffic/performance models required

Here: Comparison of 3 models & simulation (TCP Reno)

- TCP-UMass [Padhye et al., Sigcomm 98]
- TCP-Engset [Heyman et al., Sigmetrics 97]
- TCP-NBurst [modific. of Schwefel, Infocom 01]

Network Scenario I: Multiplexed ON/OFF Traffic [without wireless link]



- Scenario: N (LRD) ON/OFF sources, queueing/loss only at bottleneck router
- Models without flow-control \rightarrow known results in literature
- Here: Investigate average throughput per connection in TCP setting
- Two scenarios: (A) fast access, slow trunk, (B) vice versa

Network Scenario II: Parameter Settings

	Scenario (A)	Scenario (B)
access link speed	10 Mbit/s	128 kbit/s
trunk speed	1 Mbit/s	1.5 Mbit/s
packet size (fixed)	1000 Byte	576 Byte
average conn. size	50 kByte	200 kByte
access link speed	$\lambda_p = 1250$ packets/s	$\lambda_p = 27.78$ packets/s
trunk speed	$\nu = 125$ packets/s	$\nu = 325.52$ packets/s
average connection size	$n_p = 50$ packets	$n_p = 347.2$ packets
Buffer-Size	$B = 16$ packets	$B = 50$ packets
Delays	$\sum_i d_i = 80$ ms	$\sum_i d_i = 300$ ms
resulting minimal RTT	$R_0 = 88.8$ ms	$R_0 = 339$ ms

Observed: • Goodput per connection i : $G_i := C_i / D_i$

• Average Goodput & weighted average:

$$\bar{G} := \frac{1}{K} \sum_{i=1}^K G_i, \quad \hat{G} := \frac{1}{\sum_{i=1}^K C_i} \sum_{i=1}^K C_i \cdot G_i$$

• where $C_i := \#$ packets in conn. i , $D_i :=$ duration of conn. i

Analytic Models I: TCP-UMass [Padhye, Firoiu, Towsley, Kurose]

Approximation of throughput T in persistent TCP connection

- Derived from sender/receiver behavior
- Input Parameters
 - **b:** # packets acknowledged by ACK [here: b=1]
 - \overline{RTT} : Average Round-Trip-Time (including queuing delay)
 - W_{max} : Maximum size of congestion window [here: $W_{max}=1024$]
 - T_0 : Average time-Out interval
 - **p:** Fraction of retransmitted packets

$$T(p, \overline{RTT}) \approx \frac{1}{\overline{RTT}} \min \left(W_{max}, \frac{1}{\sqrt{\frac{2}{3}bp} + T_0/\overline{RTT} \min \left(1, 3\sqrt{\frac{3}{8}bp} \right) p(1 + 32p^2)} \right)$$

Analytic Models II: TCP-Engset [Heyman, Lakshman, Neidhardt]

Modified Processor Sharing model on connection level:

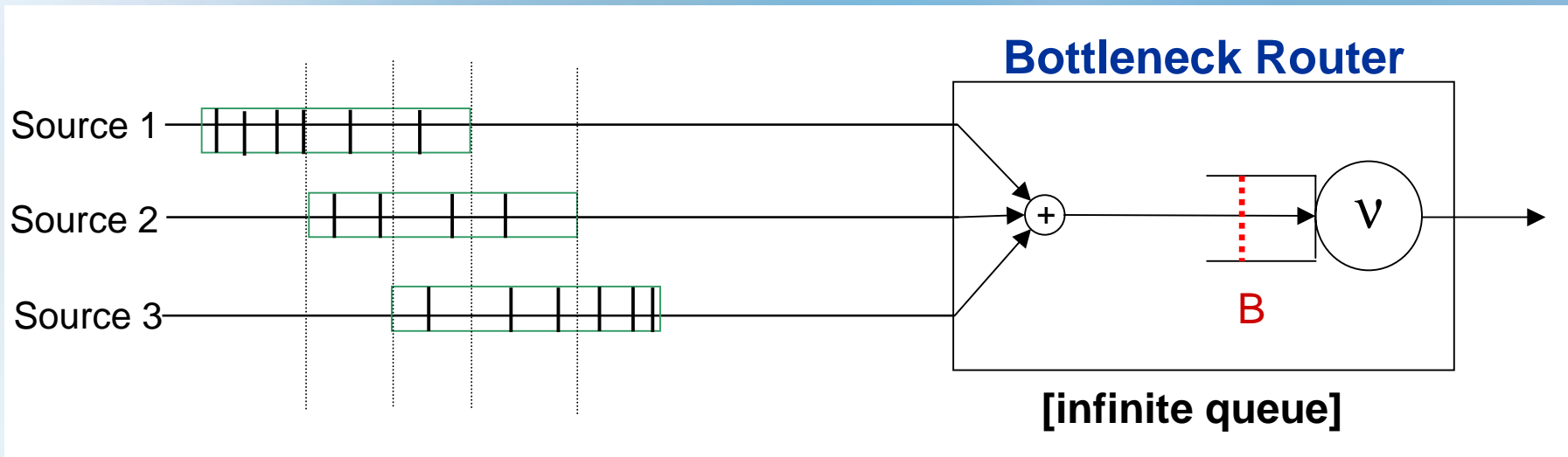
- For j active users, each obtains service at rate
 - $\sim \lambda_p$ when $j\lambda_p \leq v$
 - $\sim \delta v/j$ when $j\lambda_p > v$
- with attenuation factor $\delta \leq 1$ derived from TCP behavior

$$\delta = 1 - \frac{[(1 - p_{low})^+]^2}{2(p_{high} - p_{low})}, \quad p_{high} = \frac{B}{\nu R_0} + 1, \quad p_{low} = \frac{p_{high}}{2^{1.5}}$$

With $R_0 = \text{minimal Round-Trip Time} = \sum d_i + 1/\nu + 1/\lambda_p$

- Note:**
- # active connections can be represented by birth-death process
 - Steady-state prob. independent of distribution of connection size
(\rightarrow no impact of Long-Range Dependence by PT conn. sizes)

Analytic Models III: TCP-NBurst [Schwefel]



Cmp. TCP-Engset Packet-level extension of TCP-Engset model:

- ‚Sharing‘ of bandwidth: packet-rate λ_p at source reduced to $\delta v/j$ for j active sources when $j\lambda_p > v$
- conn. duration extended \rightarrow #packets in conn. unchanged
- Throttling only during congestion state when $Q \geq B$

Comparison (I): Applicability of Models

≡ TCP-UMass

- + Simple Formula, based on end-to-end TCP behavior
- Computation of throughput only
- ‚Derived‘ parameters required: \overline{RTT} , T_0 , p
- ? assumes persistent connections

≡ TCP-Engset

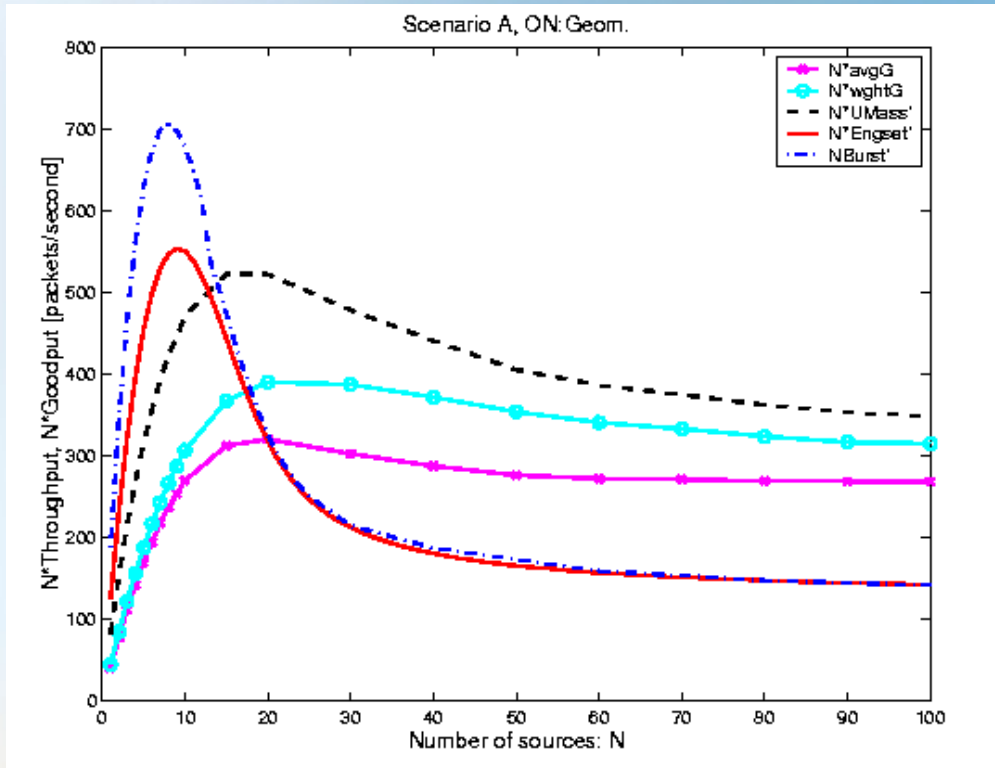
- + use of basic parameters only
- + dynamic connections (ON/OFF)
- no packet-level queue → restricted performance parameters (e.g. loss rate)
- ? Independence of connection size distribution (LRD properties)

≡ TCP-NBurst

- + packet-level queue → wide range of computable performance parameters
- computationally hard (matrix-algebraic methods), potential numerical problems

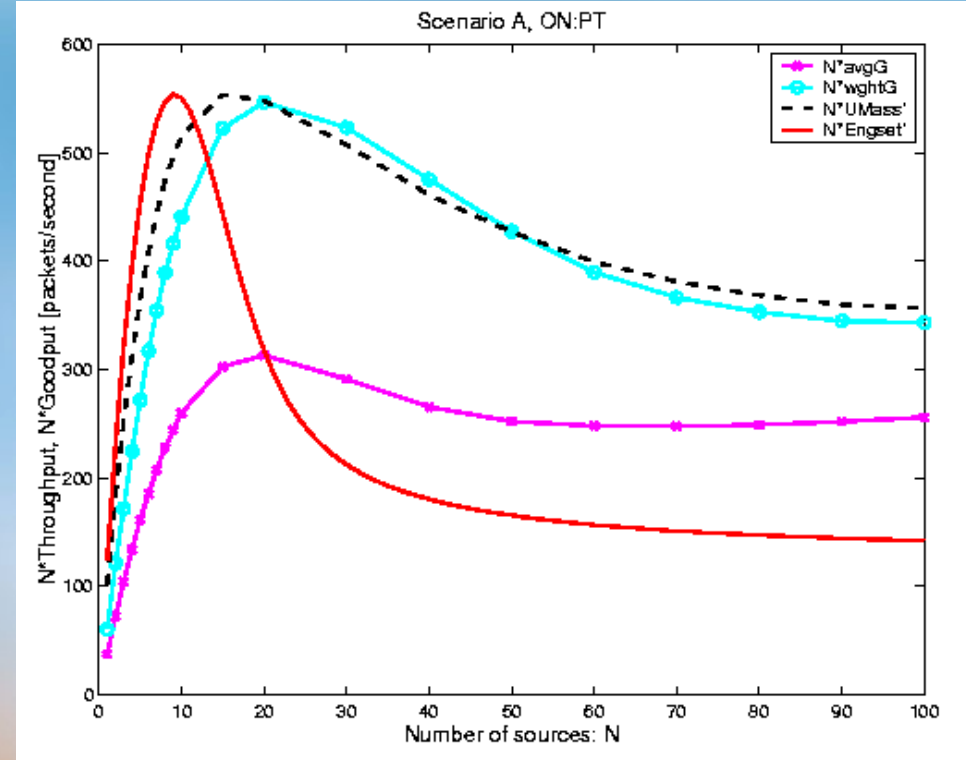
Comparison (II): Goodput per connection [Scenario (A)]

Geometric Connection Size



- TCP-UMass closest fit
- TCP-Engset and TCP-NBurst equivalent for large N (but underestimation)

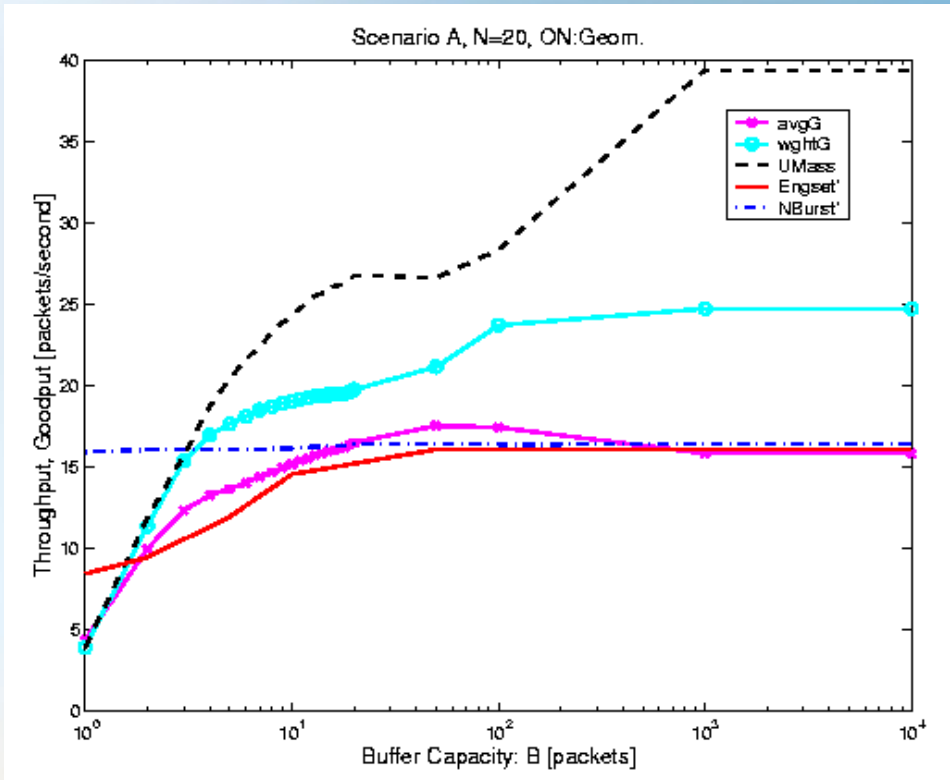
Power-Tailed Connection Size (LRD)



- Weighted Goodput \gg Plain Average
- Excellent fit of TCP-UMass to weighted average
- TCP-NBurst not computable here

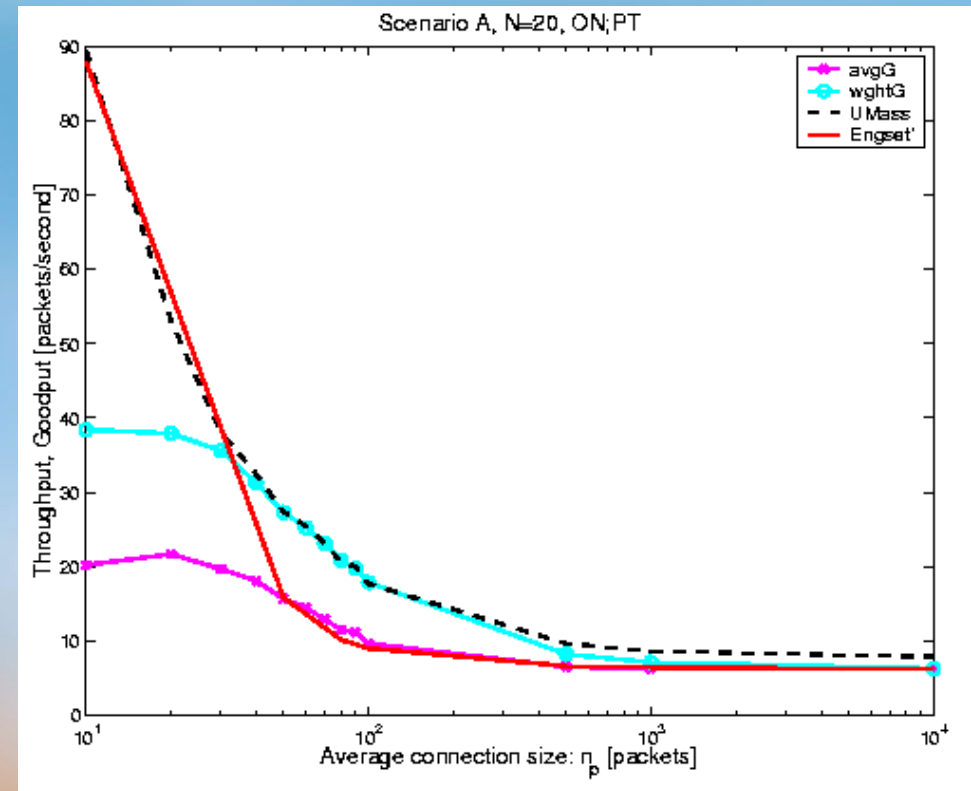
Comparison (III): Goodput per conn. [Scen. (A), N=20]

Variation of Buffer-Size



- Small buffers $B < 5$: TCP-UMass good
- Large buffers: TCP-Engset and TCP-NBurst match plain average
- TCP-Engset best match for shape of curve

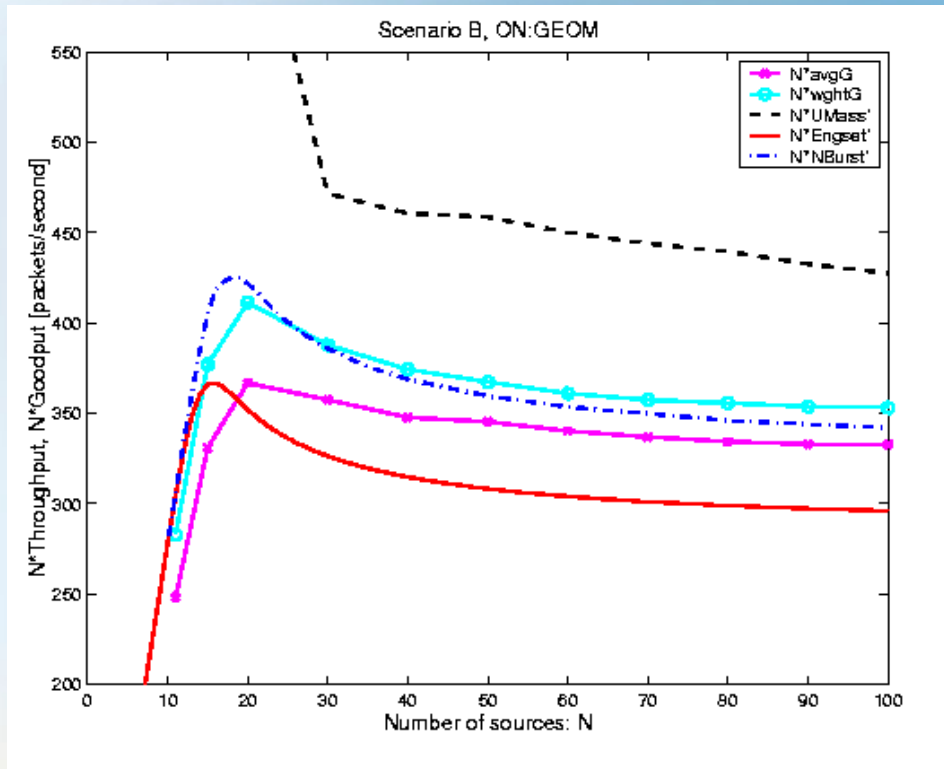
Variation of av. connection size



- All models over-estimate throughput for short connections (initial slow-start)
- Good fit for $n_p > 40$

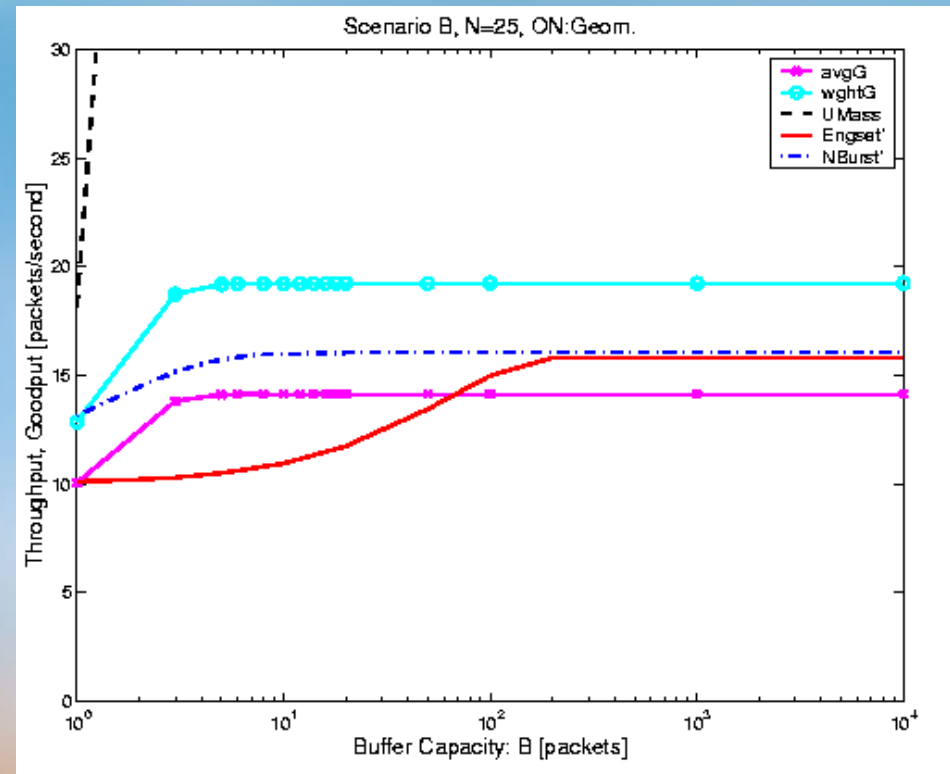
Comparison (IV): Scenario (B)

Variation of # Sources, N



- Best fit of TCP-NBurst
- TCP-UMass over-estimates largely

Variation of Buffer-size, B



- Goodput insensitive too large buffers
- TCP-UMass model useless

Summary & Outlook (Wireless scenarios)

Comparison of analytic models, TCP-UMass, TCP-Engset, and TCP-NBurst, with simulation of TCP Reno:

- ≡ TCP-UMass model requires additional (performance) parameters as input
- ≡ All models fail for short connections (initial slow-start not regarded)
- ≡ PT distributed connection sizes cause gap:
weighted average >> plain average of goodput per connection

... More experiments are required

Outlook:

- + Include properties of wireless link
 - ≡ Long delays (d_3)
 - ≡ Error model / packet-loss model for wireless link
- Necessary: Extension of analytic models for wireless links
- + Investigate impact of TCP enhancements: SACK, SNOOP